

Message authentication and hash functions

Message authentication

- Message (or document) is **authentic** if
 - It is genuine and
 - came from its alleged source.
- Message authentication is a **procedure** which verifies that received messages are authentic

Aspects of message authentication

- We would like to ensure that
 - The content of the message has not been changed;
 - The source of the message is authentic;
 - The message has not been delayed and replayed;

Message authentication techniques

- **Using conventional message encryption:**
 - if we assume that only sender and receiver share a secret key then the fact that receiver can successfully decrypt the message means the message has been encrypted by the sender
 -
- **Without message encryption**
 - The message is not encrypted, but special authentication tag is generated and appended to the message. Generation of a tag is a much more efficient procedure than encryption of the message.

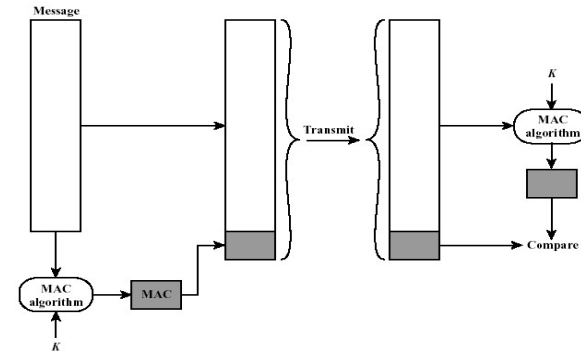
Message Authentication Code

- Let A and B share a common secret key K
- If A would like to send a message M to B , she calculates a message authentication code MAC of M using the key K :

$$MAC = F(K, M)$$

- Then A appends MAC to M and sends all this to B ;
- B applies the MAC algorithm to the received message and compares the result with the received MAC

Message authentication using MAC



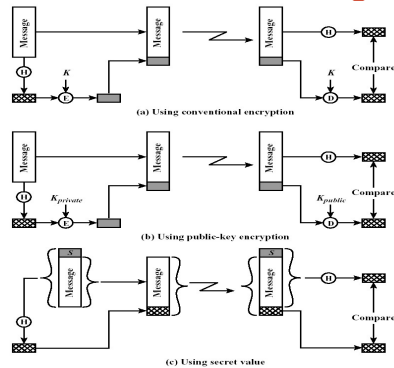
MAC algorithms

- The process of MAC generation is similar to the encryption;
- The difference is a MAC algorithm need not be reversible
→ easier to implement and less vulnerable to being broken;
- Actually, standard encryption algorithms can be used for MAC generation:
 - For example, a message may be encrypted with DES and then last 16 or 32 bits of the encrypted text may be used as MAC

One-way Hash functions

- An alternative method for the message authentication is to use one-way hash functions instead of MAC;
- The main difference is hash functions don't use a secret key:
 - $h = H(M)$;
- "One-way" in the name refers to the property of such functions: they are easy to compute, but their reverse functions are very difficult to compute.

Methods of authentication using hashes



Hash function requirements

- To be suitable for message authentication, the hash functions must have ideally the following properties:
- H can be applied to a block of data of any size;
- H produces a fixed-length output;
- $H(x)$ is easy to compute for any given x ;
- For any value h it is very difficult (infeasible) to compute x such that $H(x)=h$ (**one-way property**);
- For any given x , it is very difficult (infeasible) to find y (not equal to x) such that $H(x) = H(y)$; (**weak collision resistance**);
- It is very difficult (infeasible) to find any pair (x,y) such that $H(x) = H(y)$; (**strong collision resistance**).

Simple hash function

- Let the input be a sequence of n -bit blocks
- Then simple hash function does bit-by-bit exclusive-OR (XOR) of every block

	bit 1	bit 2	• • •	bit n
block 1	b_{11}	b_{21}		b_{n1}
block 2	b_{12}	b_{22}		b_{n2}
	•	•	•	•
	•	•	•	•
	•	•	•	•
block m	b_{1m}	b_{2m}		b_{nm}
hash code	C_1	C_2		C_n

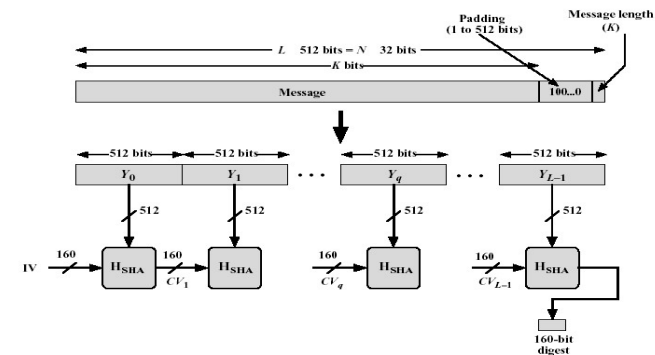
Simple hash function

- Simple hash function does not satisfy the **weak (and strong) collision property**;
- for any message M it is very easy to generate a message M_1 such that $h(M) = h(M_1)$:
 - Take arbitrary message M_2 , compute $h(M_2) = h_2$, then
 - Add additional block to M_2 , such that for the resulting M_3 we have $h(M_3) = h(M_1)$.

The SHA-1 Secure Hash Algorithm

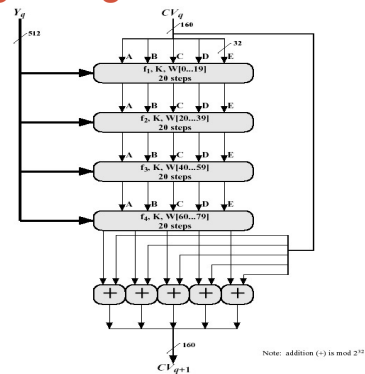
- **SHA-1 algorithm (1993-1995):**
- It has been used in the sample program illustrating password-based encryption (practical sessions);
- Takes as input a message with a maximum length
 - less than 2 to power 64 bits and produces as output a 160-bit message digest;
- The input is processed in 512-bit blocks;
- Each bit of the output is computed using all bits of the input.

SHA-1 general scheme



SHA-1 processing a single block

- The compression function;
- Includes 4 rounds with 20 steps each;
- Each round takes the current 512-bits block and 160-bit buffer value and updates the content of the buffer.



Problems and Solutions

- In 2005 a possible mathematical weakness of SHA-1 has been established:
 - ~2000 time more efficient than brute force search attack was found by Xiaoyun Wang
- Further developments: SHA-2: (SHA-224,-256,-384,-512)
- New competition for the new standard of hash functions by NIST:
 - Deadline for submissions was 31.10.2008
 - New standard SHA-3 is announced a winner on 2nd October 2012; not a replacement, but alternative for SHA-2

Recent News

- **SHAppening**, October 2015 “freestart” collision attack (by M. Stevens, P. Karpman, T. Peyrin)
- ~ US \$2000 of GPU time on EC2, est.

- **SHattered**, February 2017, full collision attack (by Google)
- ~ 6,500 years of CPU time